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# **DiCAP - An Architecture for Distributed Packet Capturing**

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Motivation  
Design and Implementation  
Evaluation  
Concluding Remarks

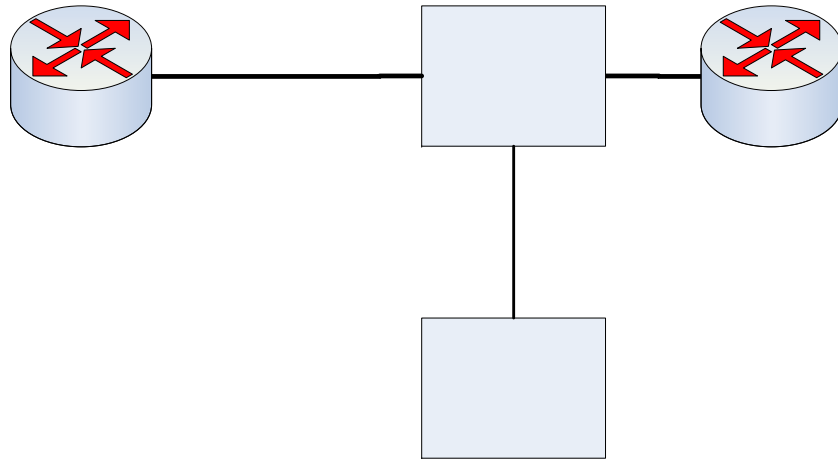


# Overview

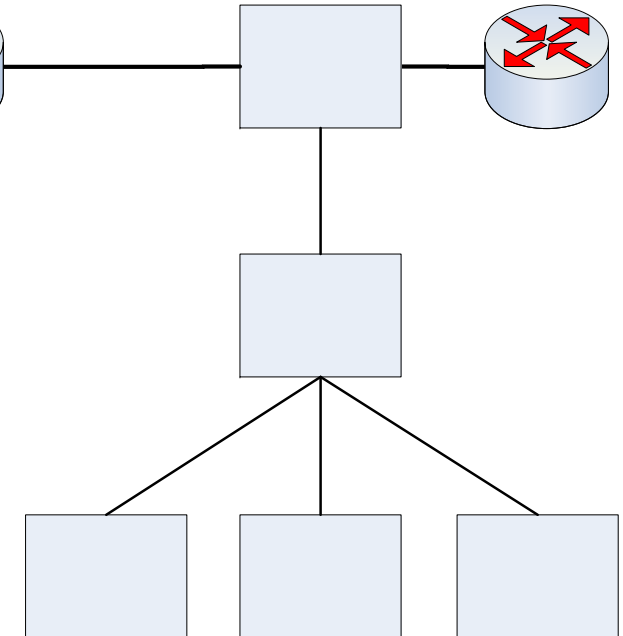
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- ❑ Most network monitoring tasks require packet inspection
- ❑ Network monitoring
  - *Live*
    - packets are inspected in real-time
    - data is dropped after inspection
    - e.g. traffic accounting, IDS
- ❑ High packet rates reduce the time available to handle a single packet
- ❑ Solutions for high packet rates:
  - packet sampling
    - decreased measurement accuracy
  - dedicated hardware
    - Expensive

# Traditional Architecture for Traffic Analysis



Mirroring  
Device



- ❑ Simple
- ❑ Easy to deploy
- ❑ Not scalable due to a single analysis box

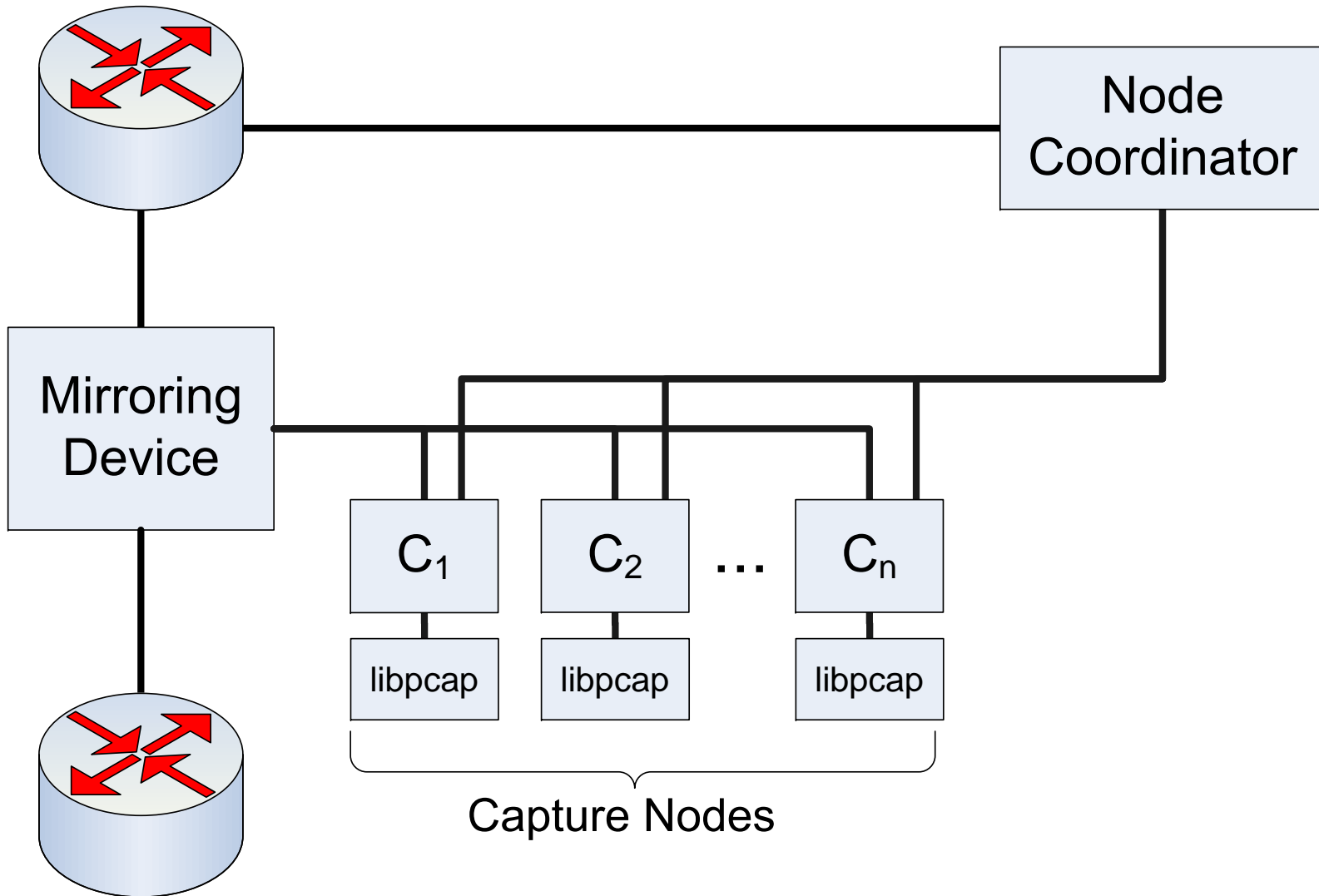
- ❑ More complex
- ❑ More scalable with respect to performance
- ❑ The single point of failure still present

# Motivation

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- Build a **distributed** architecture for IP traffic **capturing** that
  - Avoids a single point of failure
  - Avoids dedicated hardware
  - Based on off-the-shelve, inexpensive PCs
  - Allows the increase of packets that can be captured

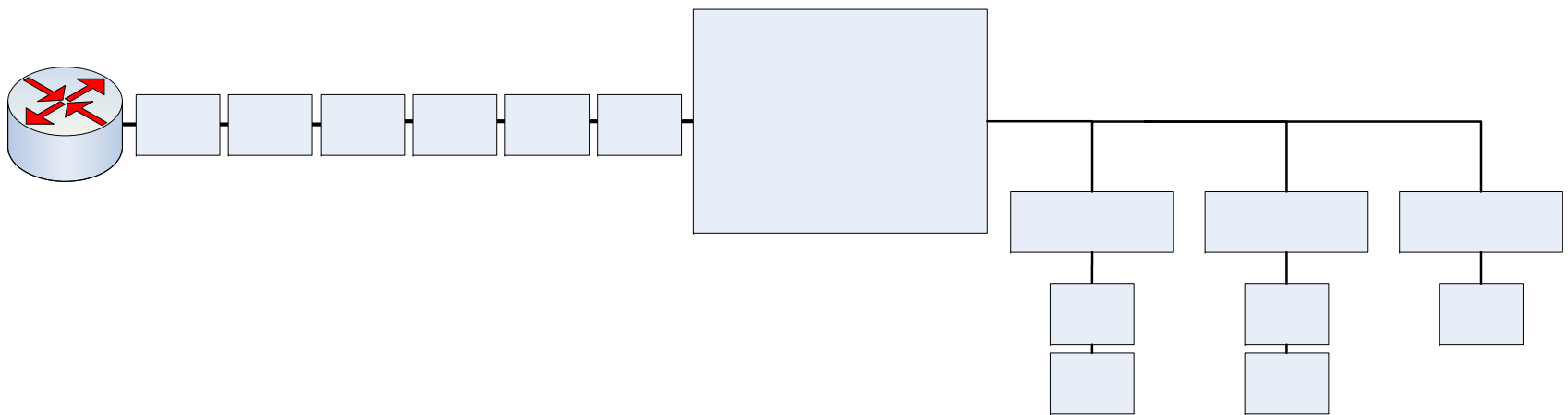
# DiCAP in Distribution Mode



# Selection Policies (1)

## □ Round robin selection

- Node coordinator introduces control packets in the mirrored traffic
- Capture nodes are logically organized in a chain
- After each packet, capture responsibility shifts to the next node in the chain
- Node coordinator configures the chain

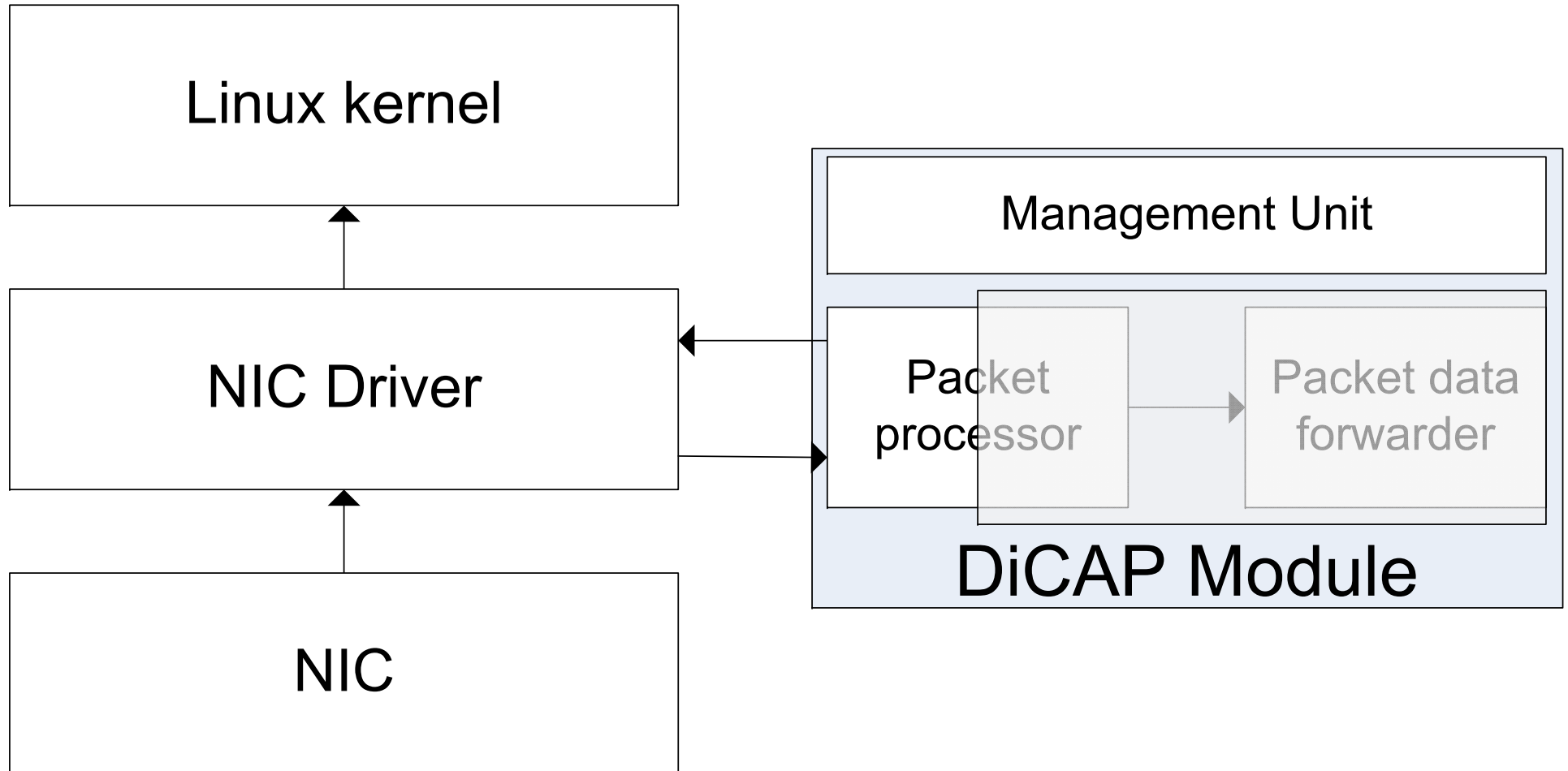


# Selection Policies (2)

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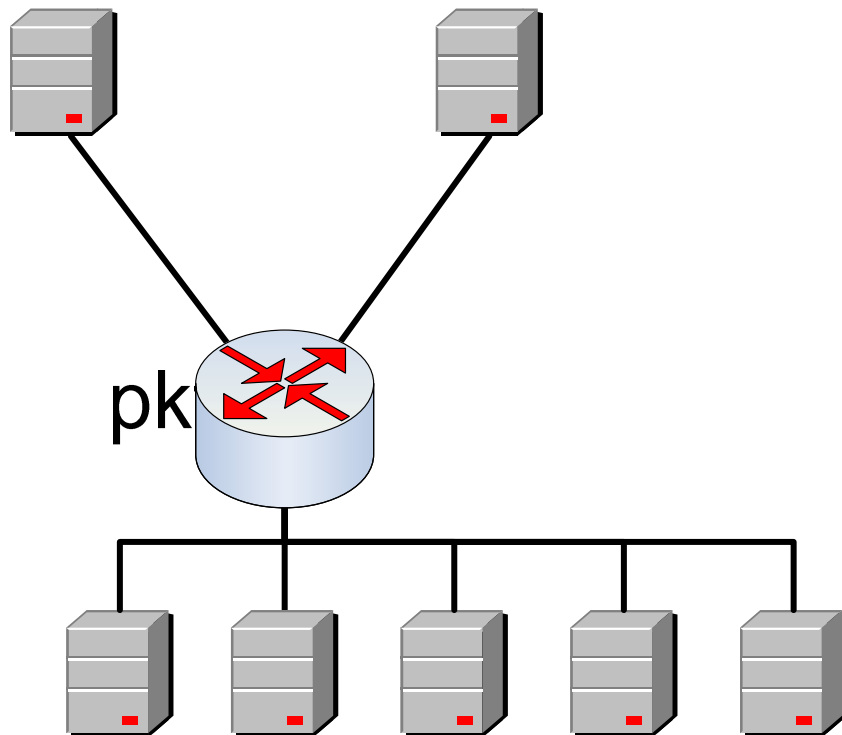
- ❑ Hash-based selection
  - A hash function is applied on packet headers
    - Current implementation uses IP identification field
  - Each node is responsible with a particular range of hash values
  - Each packet is captured by the node responsible with the respective range of hash values
  
- ❑ Advantages:
  - Easier synchronization
  - No need of control packet injection
- ❑ Disadvantages:
  - More computation needed for hash calculation

# DiCAP Implementation – Distribution Mode



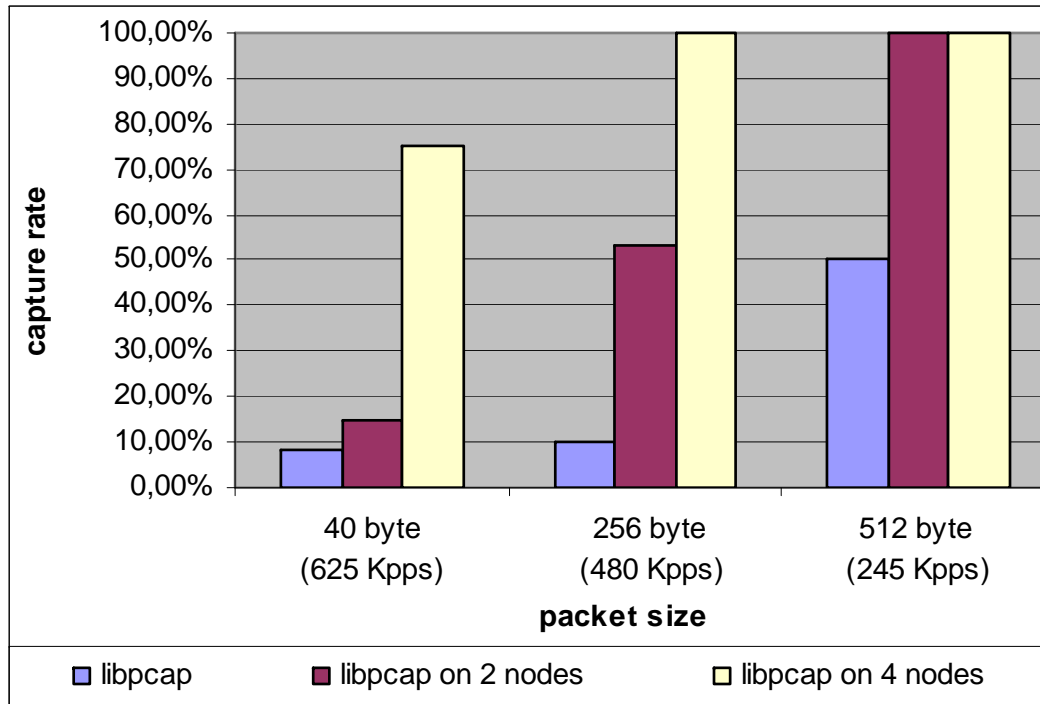


# Evaluation Testbed



- Two nodes used for traffic injection using Linux pktgen
  - One libpcap node for testing
  - Up to 4 DiCAP nodes in distribution mode
- pktgen

# Distribution Mode Evaluation



- 3 different tests were performed
  1. traditional libpcap on a single PC
  2. DiCAP in distribution mode on 2 PCs
  3. DiCAP in distribution mode on 4 PCs
- DiCAP can improve the capture performance up to 700% when 4 PCs are used in parallel

# Concluding Remarks

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- ❑ DiCAP can be used to distributedly capture packets on a high-speed link
- ❑ It may be used to allow a distributed deployment of libpcap-based applications running on common, inexpensive PCs
- ❑ It may significantly improve libpcap performance
- ❑ Very simple design, easy to implement in hardware.
- ❑ Further investigation on using other hash functions may lead to performance improvement and better load ballance