

COST 605 Meeting,  
September 24-26, 2008, Zurich, Switzerland

# PSH: A Private and Shared History-based Incentive Mechanism

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Motivation and Background  
Basic Design, PSH Algorithm  
Implementation, Evaluation



# Introduction

- P2P Systems have many advantages over centralized approaches
  - Scalability, load balancing, redundancy, no SPOF, ...
  - Popular applications: Skype, BitTorrent, Zattoo, Emule, ...



# Problem Statement

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- Problems in decentralized network
  - Free-riding, selfish behavior, malicious peers
    - Gnutella: 70% peers share no files
  - Incentives schemes necessary!
    - Tit-for-Tat (TFT), Private Shared History Incentive Mechanism (PSH)
- TFT incentive mechanism for file-sharing (BitTorrent)
  - TFT incentive mechanism works well with symmetric interest
    - A downloads from B, B downloads from A chunks from same file
    - TFT incentive mechanism uses private history (aggregated information from local observation)
- TFT incentive mechanism enables fair sharing
  - Provide unused resources now, demand resources later

# Motivation

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- TFT fails for asymmetric resource interest
  - Example DHT: Peer A stores data on B, but B does not store anything on A
  - Example IPTV: Channel switching in IPTV, new channel may require building up private history
  
- For asymmetric resource interest: Transitive TFT
  - Peer A downloads from B, B downloads from C, C downloads from A
  - Key question: How does C know about A and vice versa
    - Use shared history (observations from other peers)

# Basic Design

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- **Private Shared History Incentive Mechanism (PSH)**

- PSH combines private and shared histories

1st step: Collect information using shared history

2nd step: Evaluate and verify information using private history

# Background

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- Private history (TFT)
  - Advantages
    - Collusion resistant
    - No transaction information distribution overhead
  - Disadvantages
    - Works with symmetry of interest
- Shared history (Transitive TFT)
  - Advantages
    - Works with asymmetry of interest
  - Disadvantages
    - Not collusion resistant / not performing well
    - Transaction information distribution overhead

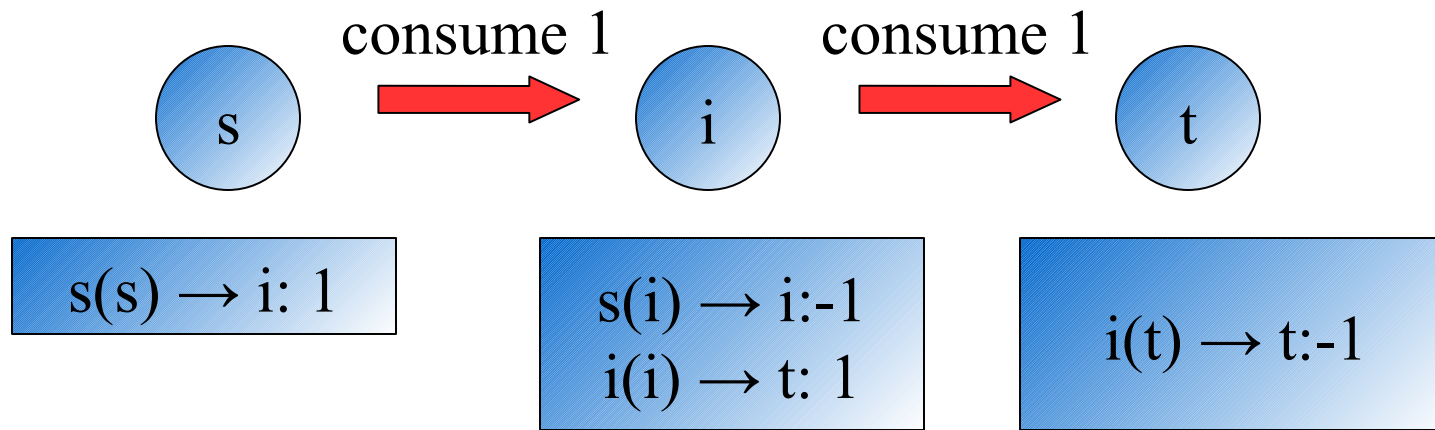
# PSH Advantages / Disadvantages

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- Private history (TFT)
  - Advantages
    - Collusion resistant
    - No transaction information distribution overhead
  - Disadvantages
    - Works with symmetry of interest
- Shared history (Transitive TFT)
  - Advantages
    - Works with asymmetry of interest
  - Disadvantages
    - Not collusion resistant / not performing well
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# PSH Algorithm (1)

- PSH algorithm in 4 steps with examples, initial credit 1
  1. Distribute latest transaction information with every request/reply
    - + no new connection required, - message size larger



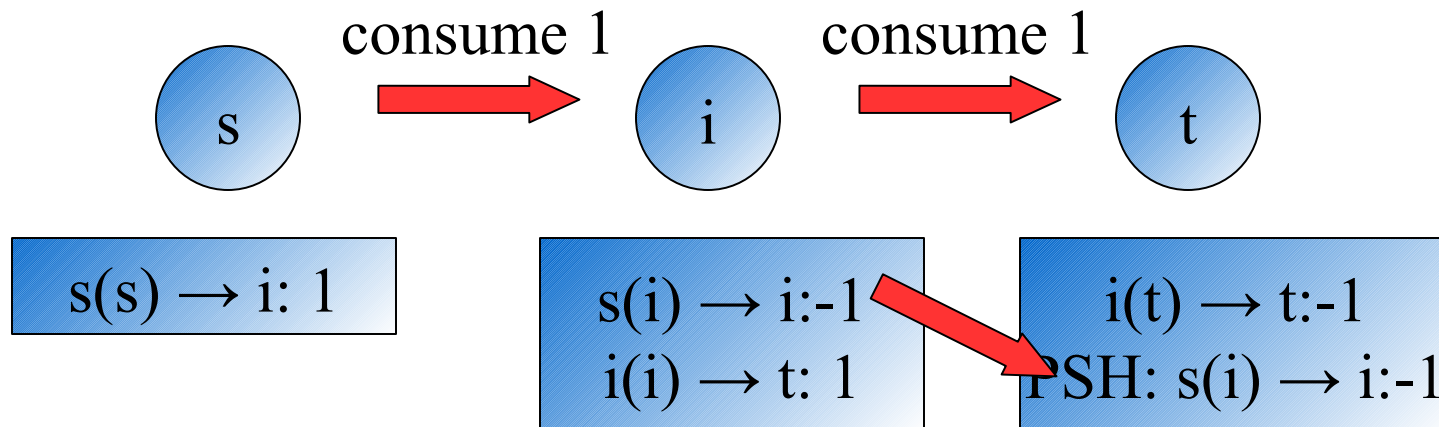
Notation  $v(w) \rightarrow x:z$

peer  $v$  consumes  $z$  resources  
from peer  $x$ , observed by peer  $w$



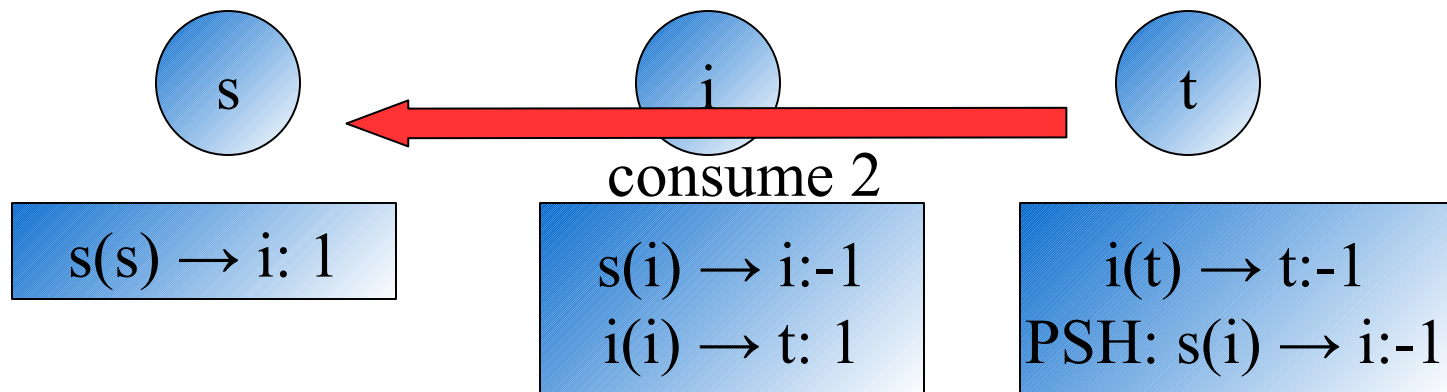
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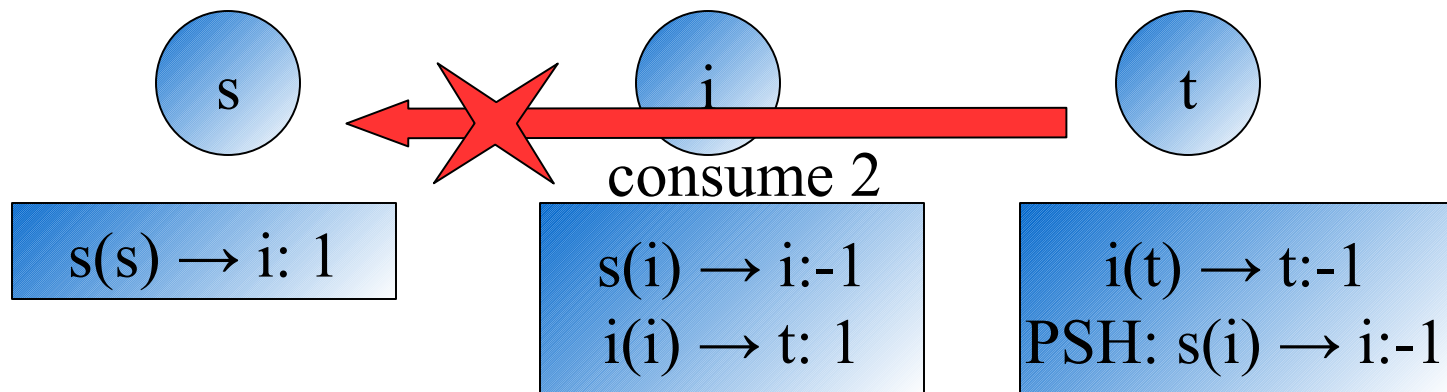
# PSH Algorithm (2)

2.If TFT fails for transaction t, path of length  $l > 2$  is searched in shared history



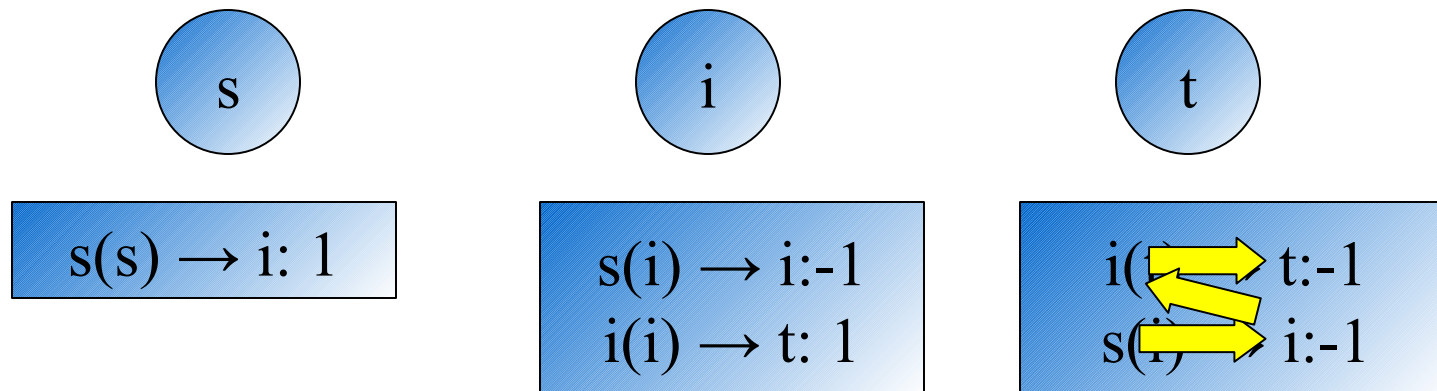
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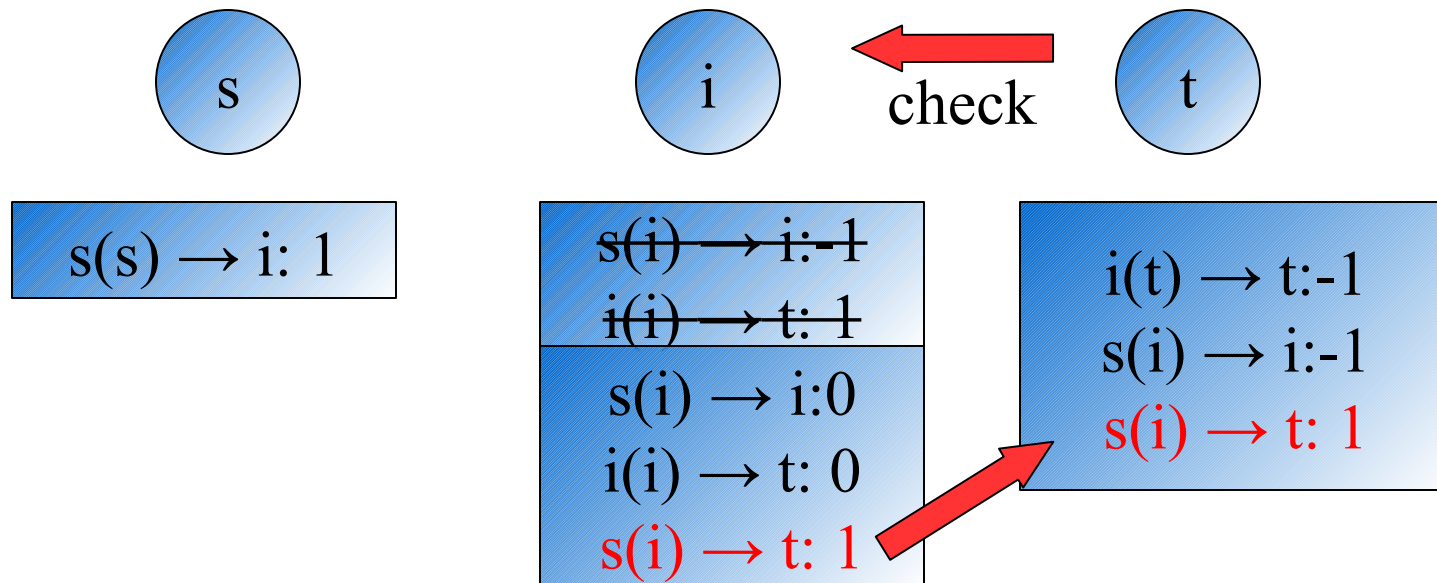
# PSH Algorithm (3)

2.If TFT fails for transaction t, path of length  $l > 2$  is searched in shared history



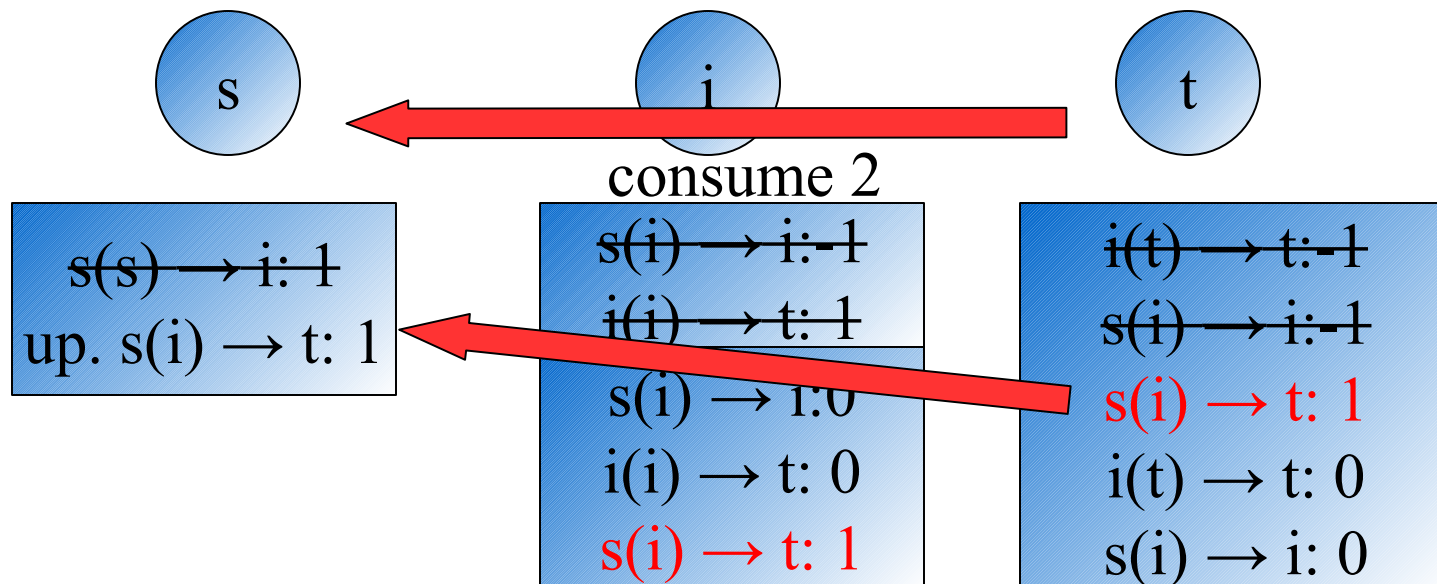
# PSH Algorithm (4)

3. If path is found, issue check c
4. Apply check c on peers involved in transaction t and execute transaction t again



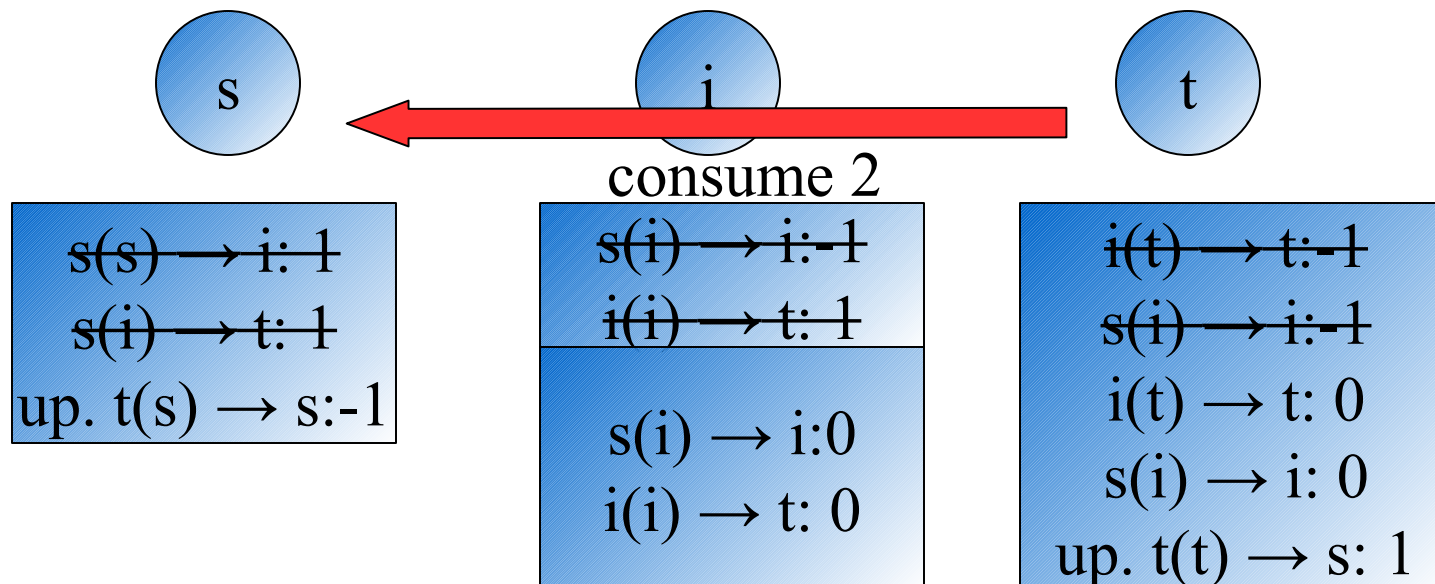
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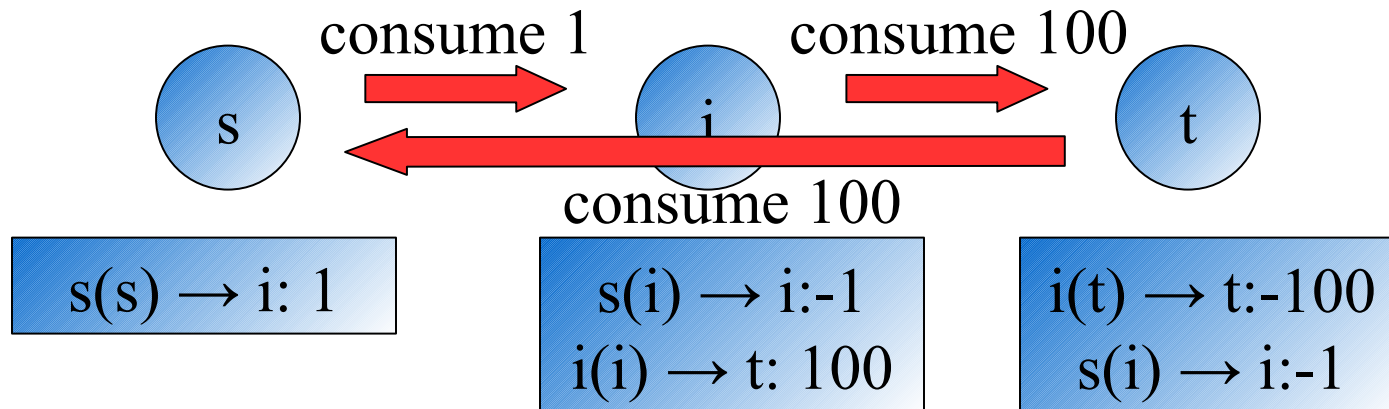
# PSH Algorithm (4)

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# PSH Algorithm – Collusion-proof

- Maximum flow for path  $s \rightarrow i \rightarrow t$  is 1



- Missing updates/wrong checks/outdated information/malicious peers
  - Peer  $s$  decides based on its private history
  - Balance on peer  $i$  remains the same
  - For peer  $t$ : amount on check may be lost



# PSH Algorithm - PSH\_r

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- ❑ PSH sends with each request latest transaction information
- ❑ PSH\_r reduces distribution overhead
  - No distribution of transaction information beforehand
  - If transaction fails, transaction information exchanged
    - + less overhead
    - + smaller messages
    - only paths of length 3 can be found reliably (one intermediate node)

# Implementation

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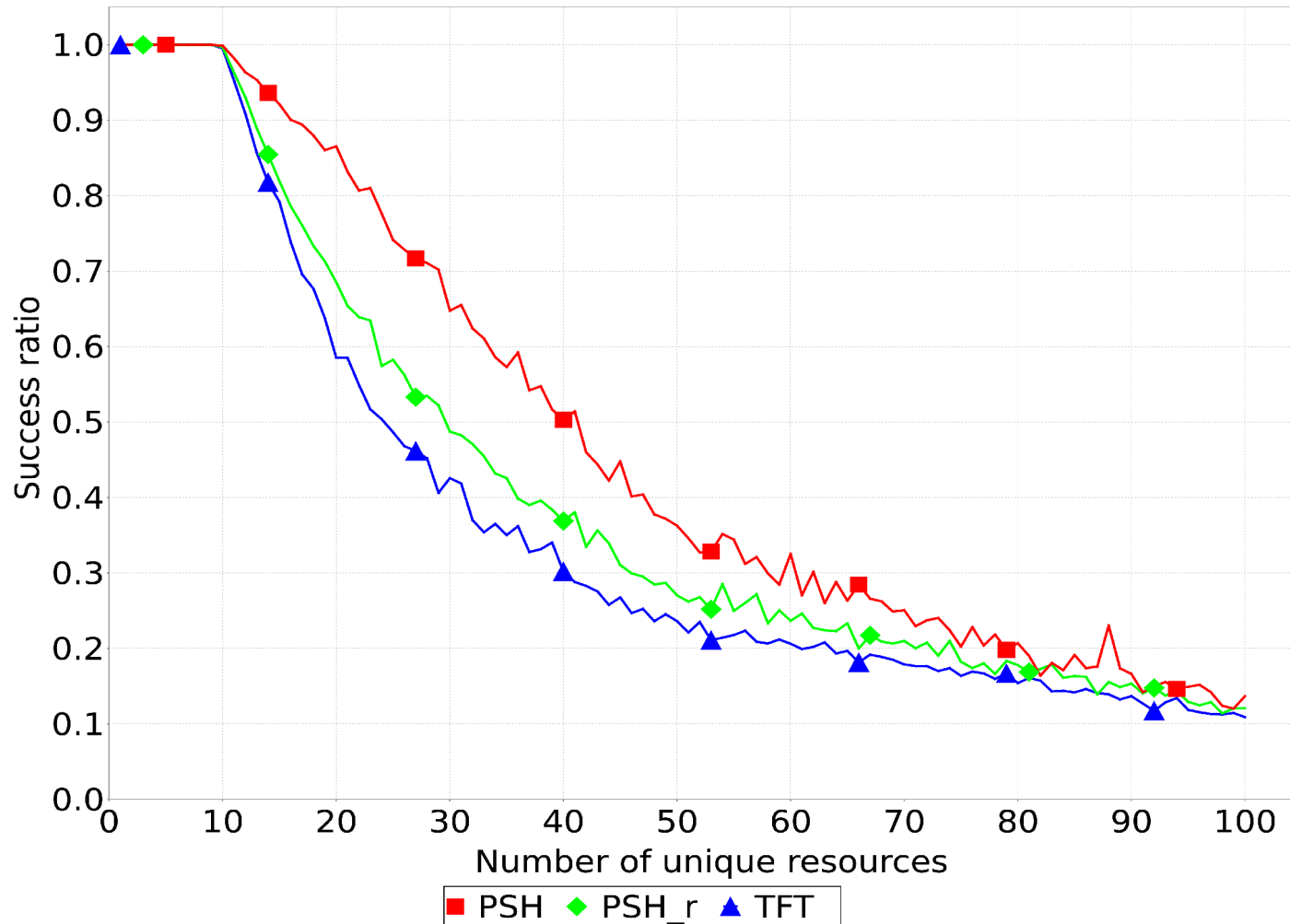
- ❑ PSH implemented in Java
- ❑ PSH uses a DHT based on Kademlia
- ❑ Evaluated with  $n$  peers,  $n=[1..N]$ ,  $N=100$
- ❑ Experiments run on one machine, operation executed sequentially
- ❑ Compared mechanisms:
  - PSH (max. 3 retries)
  - PSH\_r (no retries)
  - TFT (no retries)

# Evaluation (1)

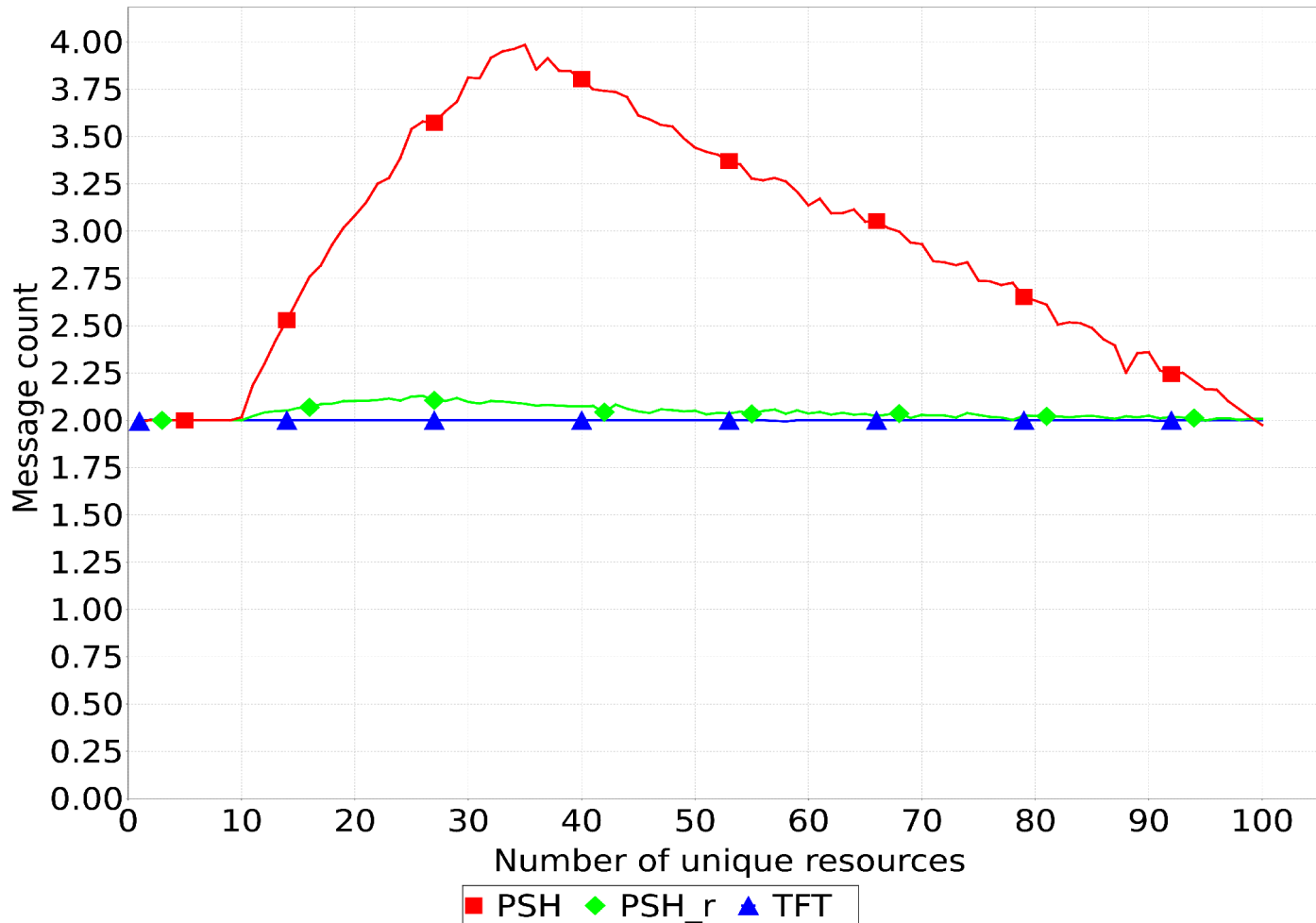
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- Parameter “number of unique resources” varied,  $p=[1..N]$ ,  $N/p$  peers share the same resource
  - $p=1$  → all peers shares the same resource (symmetric)
  - $p=100$  → all peers share different resource (asymmetric)
  - $p=50$  → 2 peers share same resource
- Plotted against:
  - Success ratio
    - With perfect knowledge success ratio is 1.0
  - Average message count
    - Number of messages per resource request/reply
  - Message size
    - Total size for all messages

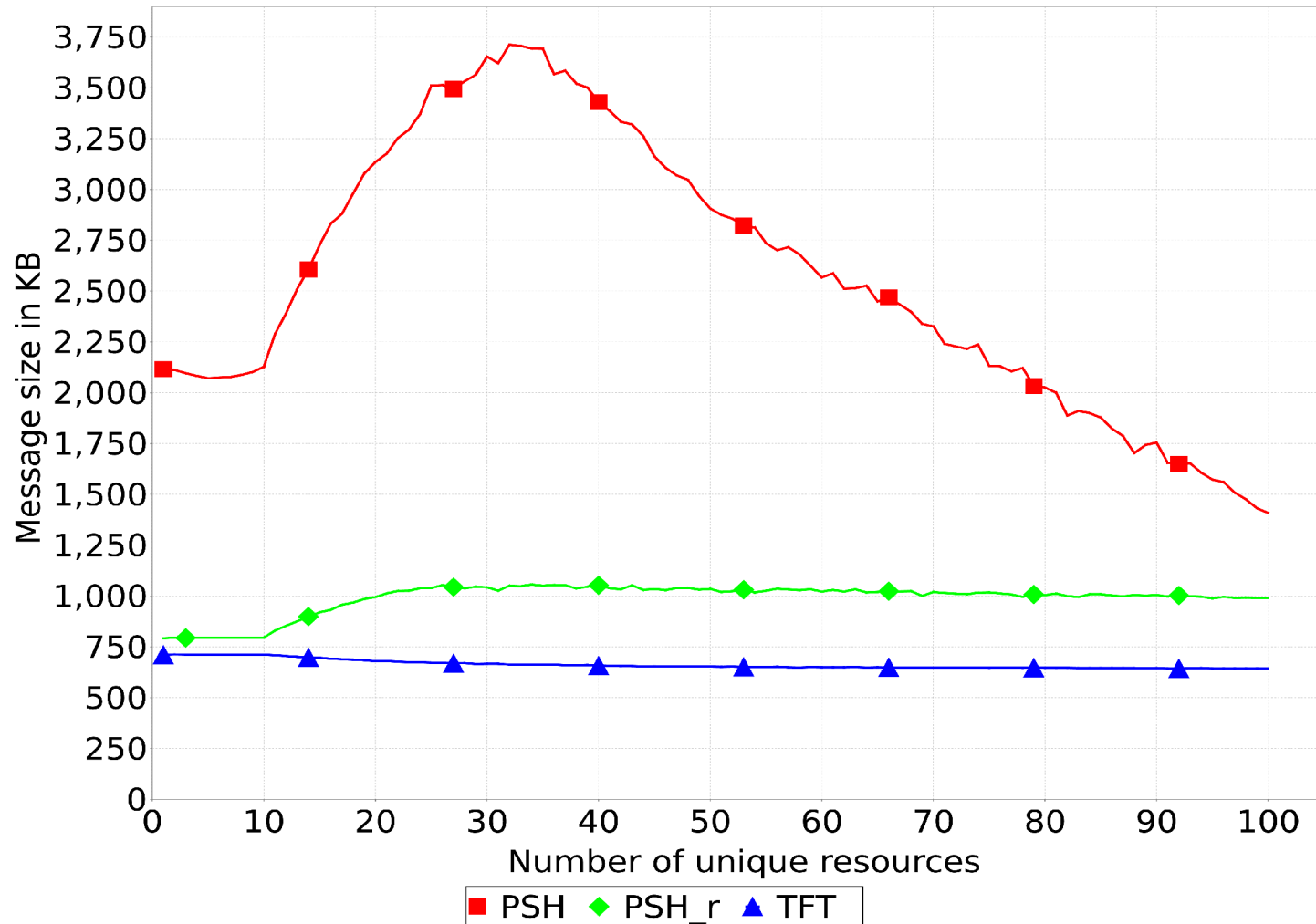
# Evaluation (2) – Success Ratio



# Evaluation (3) – Count



# Evaluation (4) – Size



# Summary and Conclusions

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- ❑ PSH combines private and shared history
- ❑ PSH works up to 73% better than TFT incentive mechanism for asymmetric resource interest
- ❑ Distribution of transaction information introduces overhead

# Future Work

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- ❑ Reduce message size: WIP PSH2
  
- ❑ Distributed testing environments
  - Deployment on EmanicsLab
  
- ❑ Possible PSH applications
  - P2P TV (channel zapping)
  - P2P Voting



# Questions?

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Thank you for your attention.

**Reference:**

T. Bocek, W. Kun, F. V. Hecht, D. Hausheer, B. Stiller: *PSH: A Private and Shared History based Incentive Mechanism*. 2nd International Conference on Autonomous Infrastructure, Management and Security Resilient Networks and Services (AIMS), July 2008