

# Sponsored Search Engines in Competition: Advertisers Behavior and Engines Optimal Ranking Strategies

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# Outline

- 1 Introduction to adword auctions & goal of the work
- 2 Model
- 3 Average prices and winning probabilities at search engines
- 4 Game between advertisers
- 5 Which mechanism to implement at the SE level?
- 6 Conclusions/Future activities

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# Introduction to adword auctions

- Search engines play a crucial role in the Internet.
- Revenue through advertising slots, usually displayed at the top or right of the search page.
- Advertisers submit bids for relevant keywords only.
- Allocation of slots thanks to adword auctions.
  - ▶ combined revenue of Yahoo! and Google in 2005: \$11 billion in 2005
  - ▶ expected to count for 40% of total advertising revenue.

The screenshot shows a Google search results page for the query "used cars usa". The search bar at the top contains the text "used cars usa" and a "Rechercher" button. Below the search bar, the results are displayed. The first result is from "Used Cars" with the URL "MoritzChevrolet.com/Used\_Cars" and the snippet "Low Prices On Used Cars - Get A Free Internet Quote!". Other results include "New & Used Cars for Sale, Auto Dealers, Car Reviews and Car ...", "Used Cars for Sale, New Cars, & Auto Buying Guide at AutoMailUsa.net", "Used Cars - Cars for Sale - Toyota - Nissan", "New Cars, Used Cars - Find Cars at AutoTrader.com", and "en stock aux Etats-Unis et au Canada - DENKER US CARS | Import New ...". On the right side of the page, there are several sponsored links, including "Used Car in Usa" and "Salvage Auto Auctions USA".

## Auction principle (single keyword, $K$ slots)

- Advertisers submit bids for specific keywords.
- Each time there is a search on that keyword:
  - ▶ advertisers are ranked and allocated slots according to a prespecified criterion:
    - ★ bid value (initially for Yahoo!)
    - ★ the revenue they will generate (more or less Google).
  - ▶ Possible payment rules:
    - ★ *Pay-Per-Impression* (PPI): advertisers charged every time their ad is displayed
    - ★ *Pay-Per-Click* (PPC): advertisers is charged only when the ad is clicked
    - ★ *Pay-Per-Transaction* (PPT): advertisers charged when a sell.
  - ▶ Amount to be paid each time?
    - ★ First Price: advertisers pay their bid
    - ★ Generalized Second Price (GSP): they pay the bid of advertiser below them in the ranking
    - ★ Vickrey-Clarke-Groves (VCG) auctions: you pay the opportunity cost that your presence introduce to all other advertisers.
- In use: PPC and GSP. **But bid-based or revenue-based ranking?**

# Goal of our work

- Almost all works deal with monopolistic search engines,
- But those engines are in competition  
⇒ What is the **best ranking strategy** given this competition?
- **Non-cooperative game** between search engines.
- Two-levels game:
  - ▶ Largest time scale: search engines choose their ranking strategy (maximizing) their revenue
  - ▶ Smallest time scale: advertisers in competition for the advertising slots (by splitting their advertisement budget over the engines)
- Game played by backward induction.
- Requires to compute the average prices and winning probabilities at search engines depending on the strategies of all actors.

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# Basic model

- Two search engines (SE), labelled 1 and 2
  - ▶ a single advertisement slot at each SE
  - ▶ considering a single keyword, with  $\lambda$ , average number of searches per unit of time
  - ▶  $\alpha$ : (fixed) market share of SE 1 ( $\alpha\lambda$  searches on SE 1).
  - ▶ Both SE apply GSP charging scheme.
- $k$  advertisers:
  - ▶ budget  $b$ , taken from cdf  $G(b)$
  - ▶ valuation per click  $v$ , taken from the cdf  $F(v)$
  - ▶ Click-Through-Rate (CTR) *separable*, as the product of the CTR of the search engine,  $q_1$  and  $q_2$  for SE 1 and 2 respectively, and of the CTR  $c_i$  of the considered Advertiser  $i$ .
  - ▶ Goal of advertiser  $i$ : split their budget  $b_i$  over the two SEs:  $\beta_i$  on SE 1 and  $1 - \beta_i$  on SE 2.
  - ▶ Remark: advertisers' interest is to bid their valuation  $v_i$  since GSP is VCG when a single slot.



## Key parameters

- $w_j(v_i)$ : probability that  $i$  wins on SE  $j$  given that her bid/valuation is  $v_i$
- $\mathbb{E}[R_j|v]$  being the average price paid on SE  $j$  having won with  $v$ .

Rates of advertisers:

$$\lambda_i^{(1)} = \min \left( \frac{\beta_i b_i}{q_1 c_i \mathbb{E}[R_1|v_i] w_1(v_i)}, \alpha \lambda \right)$$
$$\lambda_i^{(2)} = \min \left( \frac{(1 - \beta_i) b_i}{q_2 c_i \mathbb{E}[R_2|v_i] w_2(v_i)}, (1 - \alpha) \lambda \right).$$

First term: number of times Advertiser  $i$  can bid in order to reach her budget repartition, in average.

Probability that  $i$  bids on 1 and 2:

$$p_i^{(1)} = \min \left( \frac{\beta_i b}{\alpha \lambda q_1 c_i \mathbb{E}[R_1|v_i] w_1(v_i)}, 1 \right)$$
$$p_i^{(2)} = \min \left( \frac{(1 - \beta_i) b}{(1 - \alpha) \lambda q_2 c_i \mathbb{E}[R_2|v_i] w_2(v_i)}, 1 \right).$$

Revenues, for  $\beta = (\beta_1, \dots, \beta_k)$  profile of strategies of advertisers:

$$U_i(\beta) = q_1 c_i w_1(v_i) \lambda_i^{(1)} (v_i - \mathbb{E}[R_1|v_i]) q_2 c_i w_2(v_i) \lambda_i^{(2)} (v_i - \mathbb{E}[R_2|v_i]).$$

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## With bid-based ranking

If  $\pi_\ell = p_\ell$  if  $\ell$  bids,  $1 - p_\ell$  otherwise:

- Probability to win with valuation  $v$ : probability the  $k - 1$  other have valuation less than  $v$

$$w(v) = \sum_{j_1, \dots, j_{k-1} \in \{0,1\}} (F(v))^{(\sum_{l=1}^{k-1} j_l)} \left( \prod_{\ell=1}^{k-1} \pi_\ell \right)$$

*A priori* expected probability of an advertiser with valuation  $v$ .

- $\mathbb{E}[R|v]$ : expected value of the second largest bid, taking into account the probabilities that each bidder submits a bid:

$$\sum_{j_1, \dots, j_{k-1} \in \{0,1\}} \left( \prod_{\ell=1}^{k-1} \pi_\ell \right) \int_{u \leq v} \frac{(\sum_{l=1}^{k-1} j_l) u (F(u))^{(\sum_{l=1}^{k-1} j_l) - 1} f(u)}{(F(v))^{(\sum_{l=1}^{k-1} j_l)}} du.$$

Simplified expressions symmetric advertisers,  $k = 2$  and/or specific distributions (see the proceedings).

⇒ the values of  $p_i$ ,  $\mathbb{E}[R|v_i]$  and  $w(v_i)$  can be derived.

## With revenue-based ranking

- Probability to win with valuation  $v$  and CTR  $cq$ : her product  $qc_v$  larger than the  $qc_\ell V_\ell$  of each other bidding advertiser  $\ell$

$$w(v) = \sum_{j_1, \dots, j_{k-1} \in \{0,1\}} \prod_{\ell=1}^{k-1} \left( F \left( v \frac{c}{c_\ell} \right) \right)^{j_\ell} \pi_\ell.$$

- Revenue cdf of  $\ell$ :  $\mathbb{P}[qc_\ell V \leq x] = F(x/(qc_\ell))$ . Integrating the (conditional) density of the second highest revenue, we get  $cq\mathbb{E}[R|v]$ :

$$\sum_{j_1, \dots, j_{k-1} \in \{0,1\}} \left( \prod_{\ell=1}^{k-1} \pi_\ell \right) \int_{x \leq vqc} \frac{\sum_{\ell: j_\ell=1} x \frac{1}{qc_\ell} \prod_{l=1; l \neq \ell}^{k-1} (F(x/(qc_l)))^{j_l} f(x/(qc_\ell))}{\prod_{\ell=1}^{k-1} \left( F \left( v \frac{c}{c_\ell} \right) \right)^{j_\ell}} dx.$$

⇒ The values of  $p_i$ ,  $\mathbb{E}[R|v_i]$  and  $w(v_i)$  can be derived too.

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# Game between advertisers on budget repartition at each SE

We consider two SEs and two advertisers.

- The two advertisers play, trying to (selfishly) maximize their own utility/revenue.
- Equilibrium notion, **Nash equilibrium**: profile of proportion strategies  $(\beta_1^*, \beta_2^*)$  st  $\forall \beta_1, \beta_2 \in [0, 1]$ ,

$$U_1(\beta_1^*, \beta_2^*) \geq U_1(\beta_1, \beta_2^*) \text{ and } U_2(\beta_1^*, \beta_2^*) \geq U_2(\beta_1^*, \beta_2).$$

- To determine the Nash equilibria (if any), we define the *best response* of each advertiser as a function of the strategy of its opponent:

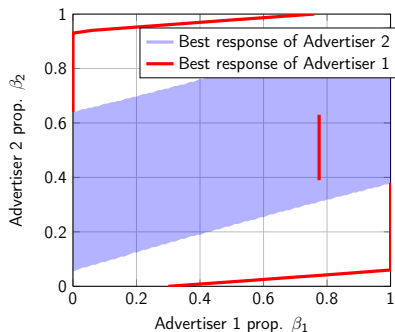
$$BR_1(\beta_2) = \arg \max_{\beta \in [0,1]} U_1(\beta, \beta_2) \text{ and}$$

$$BR_2(\beta_1) = \arg \max_{\beta \in [0,1]} U_2(\beta_1, \beta).$$

- Graphically, the set of Nash equilibria is the (possibly empty) set of intersection points of BR curves.

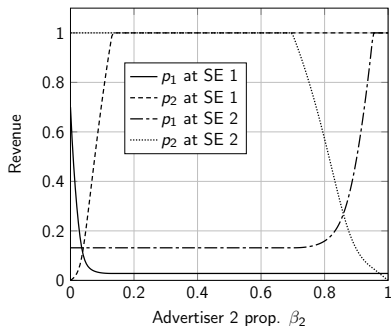
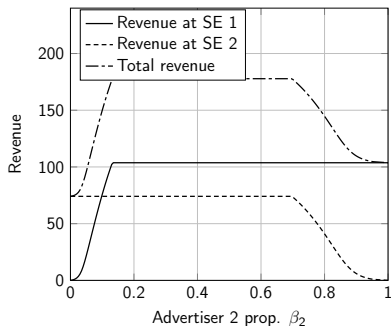
## Illustration: both SEs implement bid-based pricing

$V \equiv U[0, 20]$ ,  $\alpha = 0.6$ ,  $\lambda = 100$ ,  $q_1 = 0.5$ ,  $q_2 = 0.6$ ,  $b_1 = 5$ ,  $c_1 = 0.5$ ,  $b_2 = 20$ ,  $c_2 = 0.4$ ,  $v_1 = 10$ ,  $v_2 = 9$  and  $p_r = 0.1$



- for each fixed  $\beta_1$ , there is actually an interval for the best response of Advertiser 2 (blue)
- best response of Advertiser 1 in terms of  $\beta_2$ , we obtain the red curve
- set of Nash equilibria:  $\{0.775\} \times [0.39, 0.63]$ .

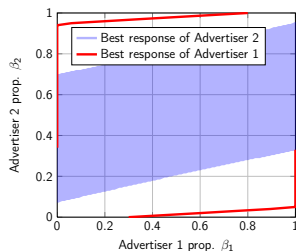
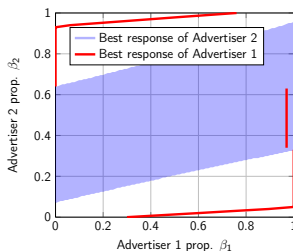
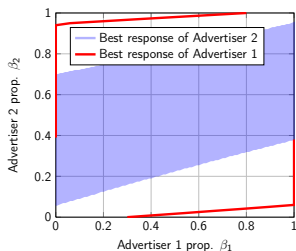
# Explanation of an interval as best response, for $\beta_1 = 0.21$



- Left: revenue indeed constant
- Right: bidding probability constant
- The probability of bidding is maximal, equal to 1, independent of the submitted budget  $\Rightarrow$  the budget is not fully spent.



### 3 other cases of ranking possibilities: bid based- revenue based (B-R), R-B and R-R



#### Nash equilibria:

- For the B-B case, all the profiles  $(\beta_1, \beta_2) \in \{0.775\} \times [0.39, 0.63]$ ;
- for the B-R case, it is  $\{0\} \times [0.39, 0.695]$ ;
- for the R-B case,  $\{0.97\} \times [0.34, 0.63]$ ;
- for the R-R case,  $\{0\} \times [0.34, 0.695]$ .

Remark: when SE 2 implements revenue-based ranking, Advertiser 1's strategy at a Nash equilibrium is to put all her budget on SE 2.

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## Game between SEs on the ranking strategy

Anticipating the advertisers' decisions, the SEs seek to maximize their revenues from advertisement.

**Another level of game:**  $(Rev_1, Rev_2)$  in terms of the rules used by SE 1 (line) and SE2 (column)

	B	R
B	(10.15, 3.62)	(1.20, <b>11.06</b> )
R	( <b>11.32</b> , 1.32)	( <b>1.50</b> , <b>11.06</b> )

- Best responses in red. Nash equilibrium: both elements in red.
- R-R: Nash equilibrium: in agreement with Yahoo!'s move to follow Google.
- For other sets of parameters such that B-R is an equilibrium.  
⇒ close look necessary for SEs!

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## Conclusions/Future activities

We have

- defined a model describing two search engines in competition for advertisers in adword auctions
- derived how advertisers should behave depending on the ranking strategies of engines
- Illustrated how engines can (competitively) play on the ranking strategy.

Future activities:

- extend our study to the situation where SEs propose more than one slot
- look more closely at Google against Yahoo!